Tutorial 8

Exercise 1: Simulate the computation of the following program for RAM that obtains a sequence (of length 1) consisting a single number 4 as an input, i.e., write down a sequence of configurations through the RAM goes during this computation:

0: $R_2 := READ()$ 1: $R_1 := 2$ 2: goto 5 3: $R_1 := R_1 * R_1$ 4: $R_2 := R_2 - 1$ 5: if $(R_2 > 0)$ goto 3 6: WRITE (R_1) 7: halt

Determine what is computed by this program, i.e., what it will produce as an output when it obtains a number n as an input.

Exercise 2: Determine what the following program for RAM computes, i.e., describe in detail its behaviour for an arbitrary input and describe what will be on the output. *Remark:* For clarity, addresses of instructions are not given explicitly, symbolic labels are used instead.

$R_4 := 4$	$R_0 := [R_1]$
$R_3 := \text{READ}()$	$R_0 := R_0 + 1$
$R_1 := R_4 + R_3$	$[R_1] := R_0$
$R_0 := 0$	goto L_2
$L_1: \ \mathbf{if} \ (R_1=R_4) \ \mathbf{goto} \ L_2$	$L_3: R_2 := 1$
$[R_1] := R_0$	$L_4: \ \mathbf{if} \ (R_2 > R_3) \ \mathbf{goto} \ L_5$
$R_1 := R_1 - 1$	$R_1 := R_4 + R_2$
goto L ₁	$R_0 := [R_1]$
$L_2: R_2 := READ()$	WRITE (R_0)
if $(R_2 \leq 0)$ goto L_3	$R_2 := R_2 + 1$
$\mathbf{if} \ (R_2 > R_3) \ \mathbf{goto} \ L_3$	goto L_4
$R_1 := R_4 + R_2$	L_5 : halt

Exercise 3: For each of the following problems, design a program for a RAM solving the given problem.

Remark: It is not necessary to deal with wrong data on input that do not correspond to specifications of inputs.

- a) INPUT: integers x, y (i.e., $x, y \in \mathbb{Z}$) OUTPUT: value x + y
- b) INPUT: integers x, y (i.e., $x, y \in \mathbb{Z}$)

OUTPUT: $\max\{x, y\}$

c) INPUT: natural number n (i.e., $n \in \mathbb{N}$) OUTPUT: sequence of numbers 1, 2, ..., n

Remark: The sequence on output will be empty for n = 0.

d) INPUT: sequence of numbers $a_1, a_2, \ldots, a_n, 0$, where $n \ge 0$ and $a_i \in \mathbb{Z} - \{0\}$ for $1 \le i \le n$ OUTPUT: $\prod_{i=1}^{n} a_i$

Remark: Notation $\prod_{i=1}^{n} a_i$ denotes the product $a_1 \cdot a_2 \cdot \ldots \cdot a_n$. For n = 0, the output will be 1.

e) INPUT: sequence of numbers $a_1, a_2, \ldots, a_n, 0$, where $n \ge 0$ and $a_i \in \mathbb{Z} - \{0\}$ for $1 \le i \le n$ OUTPUT: sequence of numbers $a_n, a_{n-1}, \ldots, a_1$

Exercise 4: Construct a program for a RAM that reads a number n from the input and writes the n-th Fibonacci number on the output. You can assume that the number n on the input is nonnegative (i.e., you don't have to consider the situation when n < 0). Recall that Fibonacci numbers F_0, F_1, F_2, \ldots are defined by the following recurrence relation:

 $F_n = \begin{cases} 0 & \text{for } n = 0 \\ 1 & \text{for } n = 1 \\ F_{n-1} + F_{n-2} & \text{for } n > 1 \end{cases}$

Exercise 5: Consider the following Algorithm 1. The input of this algorithm can be an arbitrary natural number n (i.e., values of variable n can be arbitrarily big natural numbers).

Algorithm 1:

PRINTSEQ (n): print n while n > 1 do if $n \mod 2 = 0$ then | n := n/2else [n := 3 * n + 1]print n

a) Draw the control-flow graph of this algorithm.

- b) Describe a computation performed by this algorithm when it gets number 5 as an input. Write down the sequence of configurations in this computation.
- c) How many steps are performed by the algorithm when the input is number 7? What will be the output in this case?
- d) Construct a program for RAM implementing this algorithm.

Remark: This program for RAM should read the value of value of variable n from the input.

Exercise 6: Consider Algorithm 2 described by a pseudocode.

Algorithm 2: Insertion sort

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INSERTION-SORT ():

n := \text{READ}()
for i := 0 to n - 1 do

A[i] := \text{READ}()
for j := 1 to n - 1 do

x := A[j]
i := j - 1
while i \ge 0 and A[i] > x do

A[i + 1] := A[i]
i := i - 1
A[i + 1] := x
for i := 0 to n - 1 do

WRITE (A[i])
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- a) Draw a control-flow graph representing this pseudocode.
- b) Implement this algorithm as a program for RAM.
- c) Describe how different parts of your program for RAM correspond to edges of the controlflow graph.

Exercise 7: Describe how to construct, for an arbitrary Turing machine \mathcal{M} , a program for RAM that performs the same algorithm as machine \mathcal{M} . Consider the following variants of Turing machines:

- a) A Turing machine with one tape infinite on one side
- b) A Turing machine with one tape infinite on both sides

c) A Turing machine with several tapes infinite on both sides

Remark: It is not necessary to explicitly construct these programs for RAM. It sufficient to describe informally the behaviour of these machines.

Exercise 8: Construct a program for a RAM that reads two numbers x and k from the input a writes the value of k-th bit of the number x (i.e., 0 or 1) on the output. The bits are numbered starting from 0 and 0-th bit is the least significant bit. You can assume that $x \ge 0$ and $k \ge 0$ (i.e., you don't have to consider the cases when x < 0 or k < 0).

Exercise 9: Construct a program for a RAM that reads two numbers x and y from the input (you can assume that $x \ge 0$ and $y \ge 0$) and writes their product $x \cdot y$ on the output. To make this job more difficult, you must conform to the following restrictions:

• In your program, you *must not* use arithmetic instructions for multiplication and division. However, you can use the following arithmetic instruction that implements a shift to the right by one bit:

$$R_i := rshift(R_i)$$

This instruction basically does exactly the same thing as the following intruction: $R_i := \lfloor R_i / 2 \rfloor$.

- The total number of instructions performed by your program must be polynomial with respect to the numbers of bits of numbers x and y.
- Do you have some idea how to solve this problem without instructions of the form $R_i := rshift(R_j)$, i.e., how to compute the value $x \cdot y$ on a RAM that can use only addition and subtracktion as arithmetic operations in such a way that the total number of steps is polynomial with respect the number of bits of numbers x and y?